UNCLASSIFIED

AD 664 459

ARBITRARY STATE MARKOVIAN DECISION PROCESSES

Sheldon M Ross

Stanford University Stanford, California

January 1968

Processed for ...

DEFENSE DOCUMENTATION CENTER DEFENSE SUPPLY AGENCY



U. S. DEPARTMENT OF COMMERCE / NATIONAL BUREAU OF STANDARDS / INSTITUTE FOR APPLIED TECHNOLOGY

UNCLASSIFIED

ARBITRARY STATE MARKOVIAN DECISION PROCESSES

bу

Sheldon M. Ross

TECHNICAL REPORT NO 105

January 8, 1968

Supported by the Army, Navy, Air Force, and NASA under Contract Nonr-225(53)(NR-042-002) with the Office of Naval Research

Gerald J. Lieberman, Project Director

Reproduction in Whole or in Part is Permitted for any Purpose of the United States Government

DEPARTMENT OF STATISTICS
STANFORD UNIVERSITY
STANFORD, CALIFORNIA

NONTECHNICAL SUMMARY

A Markovian Decision Process is a process which is observed at distinct time points to be in some state. After observing the state of the system an action is chosen - corresponding to the action (and the present state) a cost is incurred and the transition probabilities for the next state are determined. A policy is any rule for choosing actions. Corresponding to each policy there is an expected long run average cost per unit time. This paper is concerned with finding an optimal policy - i.e. one whose associated average cost is minimal.

For example we might have a tool which wears out with time.

The state of the system could be the length of the tool, and the possible actions could be either to replace the tool or not. Associated with each state there would be an operating cost. Thus a policy is a rule for determining when to replace the tool and an optimal one is one which minimizes the long run average cost.

In the past most of the work in this area has been done under the assumption that the state space is countable. In this paper we let the state space be arbitrary. For example, in the tool problem given above it is natural to let the state space be the continuum of possible values of the length of the tool.

This paper presents sufficient conditions for the existence of an optimal policy and for it to be of simple type. This type - called stationary deterministic - is of the form of a function mapping the state space into the action space. For example, in the tool problem

a stationary deterministic policy would replace whenever the length of the tool is in some specified set of real numbers. The method employed is to treat the average cost problem as a limit of either the discounted cost problem or the nondiscounted n-stage problem. We also show how, in a special case, the average cost problem may be reduced to a discounted cost problem.

· 大声 銀矿區

ARBITRARY STATE MARKOVIAN DECISION PROCESSES

Sheldon M. Ross

1. Introduction

We are concerned with a process which is observed at times $t=0,1,2,\ldots$ and classified into one of a possible number of states. We let χ denote the state space of the process. χ is assumed to be a Borel subset of a complete separable metric space, and we let χ be the σ -algebra of Borel subsets of χ . After each classification an action must be chosen and we let χ assumed finite, denote the set of all possible actions.

Let $\{X_t; t=0,1,2,\ldots\}$ and $\{\Delta_t; t=0,1,2,\ldots\}$ denote the sequence of states and actions; and let $S_{t-1}=(X_0,\Delta_0,\ldots X_{t-1},\Delta_{t-1})$. It is assumed that for every $x \in X$, $k \in A$ there is a known probability measure $P(\cdot \mid x,k)$ on B such that, for some version, $P\{X_{t+1} \in B \mid X_t = x, \Delta_t = k, S_{t-1}\} = P(B \mid x,k)$ for every $B \in B$, and all histories S_{t-1} . It is also assumed that for every $k \in A$, $B \in B$, $P(B \mid \cdot,k)$ is a Baire function on X.

Whenever the process is in state x and action k is chosen then a bounded (expected) cost C(x,k) - assumed, for fixed k, to be a Baire function in x- is incurred.

A policy R is a set of Baire functions $\{D_k(S_{t-1},x)\}_{k\in A}$ satisfying $D_k(S_{t-1},x)\geq 0$ for all keA, and $\sum\limits_{k\in A}D_k(S_{t-1},x)=1$ for every (S_{t-1},x) . The interpretation being: if at time t the history S_{t-1} has been observed and $X_t=x$ then action k is chosen with probability $D_k(S_{t-1},x)$. R is said to be stationary if $D_k(S_{t-1},x)=D_k(x)$ for every S_{t-1} ; R is said to be stationary deterministic if $D_k(x)$ equals 0 or 1 for all x,k.

For any policy R, let $\phi(x,R) = \limsup_{n \to \infty} \frac{1}{n} \sum_{t=0}^{n} \mathbb{E}_{R} \left[\mathbb{C}(X_{t}, \Delta_{t}) \middle| X_{0} = X \right].$ Thus $\phi(x,R)$ is the expected average cost per unit time when the process starts in state x and policy R is used.

In [4], under the assumption that χ is denumerable, a number of results dealing with the average cost criterion were proven. The method employed was to treat the average cost problem as a limit (as the discount factor approaches unity) of the discounted cost problem. In this paper we generalize some of these results to arbitrary state spaces. We also show how to treat the average cost problem as a limit of the n-stage problem. One of the advantages of this approach is that it enables us to determine, for denumerable χ , a necessary and sufficient condition for the existence of a bounded solution to a functional equation which characterizes the optimal policy.

2. Stationary Deterministic Optimal Policies

The following theorem was originally proven by Derman [2] for the special case that χ is denumerable. The following proof is new; it makes use of a technique used by Taylor [5].

Theorem 1: If there exists a bounded Baire function f(x), $x \in \chi$ and a constant g, such that

(1)
$$g + f(x) = \min_{k \in A} \{C(x,k) + \int_{\mathbb{R}} f(y) dP(y|x,k)\} \times \kappa_X$$

then there exists a stationary deterministic policy R* such that

$$g = \phi(x,R^*) = \min_{R} \phi(x,R)$$
 for all $x \in \chi$

and R^* is any policy which, for each x, prescribes an action which minimizes the right side of (1).

Proof: For any policy R

$$E_{R} \left\{ \sum_{t=1}^{n} [f(X_{t}) - E_{R}(f(X_{t}) | S_{t-1})] \right\} = 0$$

But

$$\begin{split} & E_{R}[f(X_{t}) | S_{t-1}] = \int_{y \in X} f(y) dP(y | X_{t-1}, \Delta_{t-1}) \\ & = C(X_{t-1}, \Delta_{t-1}) + \int_{y \in X} f(y) dP(y | X_{t-1}, \Delta_{t-1}) - C(X_{t-1}, \Delta_{t-1}) \\ & \geq \min_{k \in A} \{C(X_{t-1}, k) + \int_{y \in X} f(y) dP(y | X_{t-1}, k)\} - C(X_{t-1}, \Delta_{t-1}) \\ & \geq g + f(X_{t-1}) - C(X_{t-1}, \Delta_{t-1}) \end{split}$$

with equality for R^* since R^* is defined to take the minimizing action. Hence

$$0 \stackrel{\geq}{\leq} E_{R} \left\{ \sum_{t=1}^{n} [f(X_{t}) - g - f(X_{t-1}) + C(X_{t-1}, \Delta_{t-1})] \right\}$$
or
$$(2) \qquad g \stackrel{\leq}{\leq} E_{R} \frac{f(X_{n})}{n} - E_{R} \frac{f(X_{0})}{n} + E_{R} \frac{\sum_{t=1}^{n} C(X_{t-1}, \Delta_{t-1})}{n}$$

with equality for R*. Letting n ∞ and using the fact that f is bounded, we have that $g \leq \phi(R, X_0)$ with equality for R*, and for all possible values of X_0 . QED.

Remark: Note that the above proof doesn't make use of the fact that A is finite or that C(x,k) is bounded.

Let $g_n(x)$, n = 1,2,... satisfy

(3)
$$g_1(x) = \min_{k} C(x,k)$$

 $g_{n+1}(x) = \min_{k} \{C(x,k) + \int_{y \in X} g_n(y) dP(y|x,k)\}$

Note that $g_n(x) = \min_{\substack{X \in \mathbb{R} \\ X \text{ t} = 0}} \sum_{\substack{X \in \mathbb{R} \\ X \text{ t} = 0}} \mathbb{E}_{\mathbb{R}}[C(X_t, \Delta_t) | X_0 = x].$ The following corollary was proven by Derman [2] for the denumerable case.

Corollary 1: Under the conditions of theorem 1, there is a M such that

$$|g_n(x) - ng| < M$$
 for all n,x

<u>Proof:</u> Let M' be such that |f(x)| < M'. By (2) we have that $g \le 2M' + g_n(x)$. Again from (2), by letting $R = R^*$ we have that $g \ge g_n(x) - 2M'$. QED.

Fix some state - call it 0 - and let

(4)
$$f_n(x) = g_n(x) - g_n(0)$$
 all n,x

One has from (3) that

(5)
$$g_{n+1}(0) - g_n(0) + f_n(x) = \min_{k} \{C(x,k) + \int_{y \in x} f_n(y) dP(y|x,k)\}$$

We shall now determine sufficient (and in the denumerable case necessary and sufficient) conditions for the existence of a bounded Baire function f(x) and a constant g satisfying (1).

Theorem 2: If $\{f_n\}$ is a uniformly bounded equicontinuous family of functions then

(i) there exists a bounded continuous function f(x) and a constant g satisfying (1).

(ii) $\lim_{n\to\infty} (g_{n+1}(x) - g_n(x)) = g$ for all $x \in \chi$.

<u>Proof:</u> By the Ascoli Theorem there exists a subsequence $\{f_{n_k}\}$ and a continuous function f such that $f_{n_k}(x) \to f(x)$. Now $g_{n+1}(0) - g_n(0)$ is bounded (since costs are bounded) and so we can also require that $g_{n_k+1}(0) - g_{n_k}(0) + g$. Hence by (5) and the bounded convergence theorem we have that $g + f(x) = \min \{C(x,k) + \int f(y)dP(y|x,k)\}$.

For any subsequence $\{n'\}$ of $\{n\}$ there is a sub-subsequence $\{n''\}$ such that $\lim_{n\to\infty} (g_{n''+1}(0)-g_{n''}(0))$ exists. By the above this limit must be g. Thus $g=\lim_n (g_{n+1}(0)-g_n(0))$. The result follows since 0 is any arbitrary state. QED.

If \varkappa is denumerable, then $\{f_n\}$ can always be taken to be equicontinuous by considering the discrete topology. We thus have

Corollary 2: If x is denumerable, then a necessary and sufficient condition for the existence of a bounded function f(x) and constant g satisfying (1) is that there is a $M < \infty$ such that $|g_n(x) - g_n(0)| < M$ for all n, x.

Proof: Sufficiency follows from the above theorem and necessity follows from Corollary 1. QED.

For any policy R, $\beta \epsilon(0,1)$, let $\psi(x,\beta,R) = \sum_{t=0}^{\infty} \beta^t E_R[C(X_t,\Delta_t)|X_0 = x]$.

A policy R_{β} such that $\psi(x,\beta,R_{\beta}) = \min_{R} \psi(x,\beta,R)$ for all $x \in \chi$ is said to be β -optimal.

We shall need the following result given by Blackwell [1]: If A is finite, and $C(\cdot,\cdot)$ is bounded then, for each $\beta \varepsilon (0,1)$, there is a stationary deterministic policy R_{β} which is β -optimal. Furthermore $\psi(\mathbf{x},\beta,R_{\beta})$ is the unique solution to

(6)
$$\psi(x,\beta,R_{\beta}) = \min \{C(x,k) + \beta \int \psi(y,\beta,R_{\beta})dP(y|x,k)\}$$

$$k \in A \qquad y \in \chi$$

and any policy which, when in state x, selects an action which minimizes the right side of (6) is β -optimal.

Fix some state - call it 0 - and let

(7)
$$f_g(x) = \psi(x,\beta,R_g) - \psi(0,\beta,R_g)$$

then

(8)
$$g_{\beta} + f_{\beta}(x) = \min_{k} \{C(x,k) + \beta \int_{y \in X} f_{\beta}(y) dP(y|x,k)\}$$

where

$$g_{\beta} = (1-\beta) \psi (0,\beta,R_{\beta})$$

In analagous fashion to Theorem 2 we have

Theorem 3: If $\{f_{\beta}\}$ is a uniformly bounded equicontinuous family of functions then

- (i) there exists a bounded continuous function f(x) and a constant g satisfying (1).
- (ii) $(1-\beta)V_{\beta}(x) \rightarrow g$ as $\beta \rightarrow 1^-$ for all $x \in \chi$.

Proof: Same as proof of Theorem 2.

For any stationary deterministic policy R let x(R) be the action ... chooses when in state x. We say that $\lim_{n \to \infty} R_n = R$ if, for each x, there exists $N_x \le \infty$ such that $x(R_n) = x(R)$ for all $n \ge N_x$.

The following was proven in [4] for denumerable χ . The proof for arbitrary χ is identical

Theorem 4: Under the conditions of Theorem 3

- (i) for some sequence $\beta_r 1$, $R^* = \lim_{r} R_{\beta_r}$
- (ii) if $R = \lim_{\beta r} R_{\beta r}$, where $\beta_r + 1$ then R is optimal i.e.

 $\phi(x,R) = g$ for all $x \in x$.

The following two conditions were given by Taylor [5] to prove equicontinuity of $\{f_g\}$ in the special case of a replacement process:

- (a) For every $k \in A$, $C(\cdot, k)$ is continuous.
- (b) For every $x \in \chi$, $k \in A$, P(|x,k) is absolutely continuous with respect to some σ -finite measure μ on B and it possesses a density p(y|x,k) also assumed to be a Baire function in x. Furthermore, for every $x \in \chi$, $k \in A$

$$\lim_{x'\to x} \int |p(y|x,k) - p(y|x',k)|d_{\mu}(y) = 0$$

Theorem 5: If conditions (a) and (b) are satisfied then

- (i) $|f_{\beta}(x)| < M$ for all $x, \beta = \{f_{\beta}\}$ is equicontinuous
- (ii) $|f_n(x)| < M$ for all $x, n \Rightarrow \{f_n\}$ is equicontinuous

Proof: Follows directly from (5) and (8) and conditions (a), (b).

A sufficient condition for the uniform boundedness of $\{f_{\beta}\}$ is given in [4].

3. Reduction of Average Cost Case to Discounted Cost Case

We shall need the following assumption

Assumption (I): There is a state - call it 0 - and $\alpha \ge 0$, such that $P\{X_{t+1} = 0 \mid X_t = x, \Delta_t = k\} = \alpha \text{ for all } x \in X, k \in A.$

For any process satisfying the above Assumption consider a new process with identical state and action spaces, with identical costs, but with transition probabilities now given by

$$P'(B|x,k) = \begin{cases} \frac{P(B|x,k)}{1-\alpha} & \text{for } 0 \notin B \\ \\ \frac{P(B|x,k)-\alpha}{1-\alpha} & \text{for } 0 \in B \end{cases}$$

Let $\psi'(x,\beta,R)$ be the total expected β -discounted cost, and let R'_{β} be the β -optimal policy, all with respect to the new process. Letting $f'(x) = \psi'(x,1-\alpha,R'_{1-\alpha}) - \psi'(0,1-\alpha,R'_{1-\alpha})$ we have by (8) that

(9)
$$\alpha \psi'(0,1-\alpha,R) + f'(x) = \min_{k} \{C(x,k) + (1-\alpha) \int_{0}^{\infty} f'(y) dP'(y|x,k)\}$$

$$= \min_{k} \{C(x,k) + \int_{0}^{\infty} f'(y) dP(y|x,k)\}$$

And thus the conditions of Theorem 1 are satisfied. It follows that $\mathbf{g} = \alpha \psi'(0, 1-\alpha, \mathbf{R}_{1-\alpha}')$ and the optimal average-cost policy is the one which selects the actions which minimize the right side of (9). But it is easily seen that $\mathbf{R}_{1-\alpha}'$ does exactly this. Hence the optimal average cost policy is precisely the 1- α -optimal policy with respect to the new process; and the optimal expected average cost per unit time is $\alpha \psi'(0, 1-\alpha, \mathbf{R}_{1-\alpha}')$.

The above result was proven in [4] for the denumerable case by showing that $\phi(x,R) = \alpha \psi'(0,1-\alpha,R)$ for any stationary policy R.

This result also holds for arbitrary χ . However this in itself does not show that $R_{1-\alpha}^{'}$ is optimal. (It does in the denumerable case because it can be shown that Assumption (I) implies that $\{f_{\beta}\}$ is uniformly bounded and thus by Theorem 3 there exists a stationary deterministic policy which is optimal.)

4. Concluding Remarks

Results given in [4] which dealt with ε -optimal policies and replacement processes (Sections 3 and 4) carry over to the more general spaces χ considered here. The proofs are identical (with integrals replacing sums in the obvious places).

REFERENCES

- [1] Blackwell, David (1965), Discounted Dynamic Programming,

 Annals of Mathematical Statistics, 36, 226-235.
- [2] Derman, Cyrus (1966), Denumerable State Markovian Decision
 Processes Average Cost Criterion, Annals of Mathematical
 Statistics, 37, 1545-1554.
- [3] Derman, Cyrus and Lieberman, Gerald J. (1966), A Markovian

 Decision Model for a Joint Replacement and Stocking Problem,

 Management Science, 13, 609-617.
- [4] Ross, Sheldon M. (1967), Non-Discounted Denumerable Markovian

 Decision Models, To appear in Annals of Mathematical Statistics.
- [5] Taylor, Howard (1965), Markovian Sequential Replacement Processes,

 Annals of Mathematical Statistics, 36, 1677-1694.

UNCLASSIFIED

	~		
Security	•	 11108	HAR
DACHILL	~		

DOCUMENT CO	NTROL DATA - Re		the everell report is classified)	
1. ORIG NATING ACTIVITY (Corporate author) Stanford University	•	24 REPO	AT SECURITY CLASSIFICATION INCLESSIFIED	
Department of Statistics		2 5 6 ROU	•	
Stanford, California				
3. REPORT TITLE				
Arbitrary State Markovian Decision	Processes			
4. DISCRIPTIVE NOTES (Type of report and inclusive dates) Technical Report				
S. AUTHOR(5) (Last name, Bret name, Initial)				
Sheldon M. Ross				
6. REPORT DATE	74- TOTAL NO. OF	PAGES	75. NO. OF REPS	
January 8, 1968	12		5	
Se. CONTRACT OR GRANT NO.	Se. ORIGINATOR'S R	EPORT NUM	BER(S)	
Nonr-225(53) PROJECT NO.	105			
NR-042-002				
c.	SE OTHER REPORT	NO(S) (Any	other numbers that may be essigned	
	ante reporti			
d. 10 A VAIL ABILITY/LIMITATION NOTICES	1			
TO A VA IL ABILITY/LIMITATION NOTICES				
Distribution of this document is unlim	ited			
11. SUPPLEMENTARY NOTES	12 SPONSORING MILITARY ACTIVITY Logistics and Mathematical Statistics Br. Office of Naval Research Washington, D.C. 20360			
13. ABSTRACT				
Arbitrary state, finite action Mawith respect to the (long-run) average treated both as a limiting case of the a limit of the n-stage problem. Suffice existence of an optimal rule and for in	cost criterion discounted cost cient condition	t problems are gi	oroblem is em and also as lven for the	

DD .5084 1473

UNCLASSIFIED
Security Classification

Security	O1 11	C1 A1 -

14	LIN	LINKA		LINK B		LINK C	
KEY WORDS		WT	HOLE	WT	ROLE	WT	
Markovian Decision Process					<u> </u>		
Arbitrary State Space							
Stationary Deterministic Optimal Rule							

INSTRUCTIONS

- 1. ORIGINATING ACTIVITY: Enter the name and address of the contractor, subcontractor, grantee, Department of Defense activity or other organization (corporate author) issuing the report.
- 2a. REPORT SECURTY CLASSIFICATION: Enter the overall security classification of the report. Indicate whether "Restricted Data" is included. Marking is to be in accordance with appropriate security regulations.
- 2b. GROUP: Automatic downgrading is specified in DoD Directive 5200-10 and Armed Forces Industrial Manual. Enter the group number. Also, when applicable, show that optional markings have been used for Group 3 and Group 4 as authorized.
- 3. REPORT TITLE: Enter the complete report title in all capital letters. Titles in all cases should be unclassified. If a meaningful title cannot be selected without classification, show title classification in all capitals in parenthesis immediately following the title.
- 4. DESCRIPTIVE NOTES: If appropriate, enter the type of report, e.g., interim, progress, summary, annual, or final. Give the inclusive dates when a specific reporting period is covered.
- 5. AUTHOR(S): Enter the name(s) of author(s) as shown on or in the report. Enter last name, first name, middle initial. If military, show rank and branch of service. The name of the principal author is an absolute minimum requirement.
- 6. REPORT DATE: Enter the date of the report as day, month, year, or month, year. If more than one date appears on the report, use date of publication.
- 7a. TOTAL NUMBER OF PAGES: The total page count should follow normal pagination procedures, i.e., enter the number of pages containing information.
- 7b. NUMBER OF REFERENCES: Enter the total number of references cited in the report.
- 8a. CONTRACT OR GRANT NUMBER: If appropriate, enter the applicable number of the contract or grant under which the report was written.
- 86, 8c, & 8d. PROJECT NUMBER: Enter the appropriate military department identification, such as project number, subproject number, system numbers, task number, etc.
- 9a. ORIGINATOR'S REPORT NUMBER(S): Enter the official report number by which the document will be identified and controlled by the originating activity. This number must be unique to this report.
- 9b. OTHER REPORT NUMBER(S): If the report has been assigned any other report numbers (either by the originator or by the sponsor), also enter this number(s).
- 10. AVAILABILITY/LIMITATION NOTICES: Enter any limitations on further dissemination of the report, other than those

imposed by security classification, using standard statements such as:

- (1) "Qualified requesters may obtain copies of this report from DDC."
- (2) "Foreign announcement and dissemination of this report by DDC is not authorized."
- (3) "U. S. Government agencies may obtain copies of this report directly from DDC. Other qualified DDC users shall request through
- (4) "U. S. military agencies may obtain copies of this report directly from DDC. Other qualified users shall request through
- (5) "All distribution of this report is controlled. Qualified DDC users shall request through

If the report has been furnished to the Office of Technical Services, Department of Commerce, for sale to the public, indicate this fact and enter the price, if known.

- 11. SUPPLEMENTARY NOTES: Use for additional explanatory notes.
- 12. SPONSORING MILITARY ACTIVITY: Enter the name of the departmental project office or laboratory sponsoring (paying for) the research and development. Include address.
- 13. ABSTRACT: Enter an abstract giving a brief and factual summery of the document indicative of the report, even though it may also appear elsewhere in the body of the technical report. If additional space is required, a continuation sheet shell be attached.

It is highly desirable that the abstract of classified reports be unclassified. Each paragraph of the abstract shall end with an indication of the military security classification of the information in the paragraph, represented as (TS), (S), (C), or (U)

There is no limitation on the length of the abstract. How ever, the suggested length is from 150 to 225 words.

14. KEY WORDS: Key words are technically meaningful terms or short phrases that characterize a report and may be used as index entries for cataloging the report. Key words must be selected so that no security classification is required. Identifiers, such as equipment model designation, trade name, military project code name, geographic location, may be used as key words but will be followed by an indication of technical context. The assignment of links, rales, and weights is optional.